

## Research paper

**A Hybrid Mixed-Integer Optimization and Clustering Approach to Selective Collection Services Problem of Domestic Solid Waste***Un Enfoque Híbrido de Agrupamiento y Optimización Entera Mixta para el Problema de Servicios de Recolección Selectiva de Residuos Sólidos Domésticos***Johana Andrea Patiño Chirva<sup>1</sup> , Yesica Xiomara Daza Cruz<sup>1</sup> , Eduyn Ramiro Lopez-Santana<sup>1</sup>**<sup>1</sup>Universidad Distrital Francisco José de Caldas- Facultad de Ingeniería. email: erlopezs@udistrital.edu.co

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**Abstract**

**Context:** Waste generation is causing profound negative impacts on our natural environment. Because of that, processes related to waste collection, transportation, transformation and final disposal have increased its importance and major efficiencies are desirable. We propose a Mixed Integer Programming model and clustering approach for waste collection and transportation process.

**Method:** An optimization model, inspired on Bogotá context is proposed, to maximize the amount of waste collected, considering real-life aspects of this activity in the city. For large instances in which there is a big computational cost, we proposed an alternative solution of two stages, firstly a clustering step and then a routing step.

**Results:** In small instances of up to 1453 collection sites grouped in 13 blocks and 51 corners, the model result in an overall collection covering of 100%. For large instances, there are variations between the results of each clustering method. The results suggests that the sweep algorithm is better to clustering the collection sites.

**Conclusions:** Our proposed model is able to find a solution the waste collection problem in Bogota case considering the vehicle capacity, maximum workday duration and the planning horizon of two days according with the collection process in Bogota. We test three clustering methods in order to group the collection sites and to reduce the complexity of the problem, and then we solve the model using a commercial solver. For the small instances, our model run very fast but in the large in-stances the computational time was increased. Future work will focus in the validation and search of solution methods improving the performance with the proposed model.

**Keywords:** clustering, optimization model, routing, scheduling, waste collection services.

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### Resumen

**Contexto:** La generación de residuos está causando profundos y negativos impactos en nuestro ambiente. Por esto, procesos relacionados con la recolección, el transporte, la transformación y la disposición final de residuos han ganado importancia y se propende a su eficiencia. Un modelo de programación entera mixta y de agrupación es propuesto para los procesos de recolección y transporte de residuos.

**Método:** Se propone un modelo de optimización, inspirado en el contexto de Bogotá, cuyo objetivo es maximizar la cantidad de residuos recolectados teniendo en cuenta las características reales de esta actividad en la ciudad. Para instancias grandes en las que el costo computacional es muy alto, se proponen una alternativa de solución de dos etapas, clusterizar primero y rutear después.

**Resultados:** En pequeñas instancias de hasta 1453 puntos de recolección agrupados en 13 bloques y 51 esquinas, nuestro modelo logró el cubrimiento total de la recolección. Para grandes instancias, existen variaciones entre los resultados de cada método de agrupación.

**Conclusiones:** Nuestro modelo propuesto es capaz de encontrar una solución al problema de recolección de residuos en caso de Bogotá considerando la capacidad de los vehículos, la duración máxima de una jornada de trabajo y el horizonte de planificación de dos días de acuerdo con el proceso de recolección de Bogotá. Probamos tres métodos de agrupación para agrupar los sitios de recolección y para reducir la complejidad del problema, y luego se resuelve el modelo usando un paquete comercial. Para instancias pequeñas, nuestro modelo es rápido, pero en grandes instancias se aumenta el tiempo de cálculo. Los trabajos futuros se centrarán en la búsqueda y validación de métodos con el propósito de encontrar aquel que tenga un mejor desempeño con el modelo propuesto.

**Palabras clave:** clúster, modelo de optimización, ruteo, programación de tareas, recolección de residuos.

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## 1. Introduction

The transformation of cities and the increased of urban population have had several environmental impacts, one of them concern directly with consumption habits, talking about the waste generation and the practices around its last disposition. In this context, the government of Bogotá has done some programs in order to give better destination for waste. These programs not only involve environmental aspects but also there is a social point which is taking into account, the roll of waste pickers. The purpose is to have better practices for waste final disposal by its separation at source and also the participation of recyclers in its returning into the productive cycle. To make this possible waste pickers are being organized into legal organization with better conditions than they had before.

In the progress of recycler inclusion, there is a concern that has been considered, the fact that they have not any technical criterion in the selection of the route that they usually do, so it cannot be guaranteed 100% of recyclable solid waste return which is one of the principal objectives of Bogotá management programs. For this reason, we propose an optimization model inspired on Bogotá context to maximize the amount of waste recollected, considering the real features of this activity in the city.

Waste management includes the process of collection, transport, processing, recycling and final disposition. Within this process there have been taking into consideration several factors like environment, society, legislation, economic, technology and politic in order to have a framework to make decisions such as the opening of recycling centers, final disposition, acquisition of trucks and the development of efficient routes, among others.

Problems of waste collection have been studied for general materials as solid waste [1], urban solid waste [2], commercial waste [3], rural solid waste [4], and specific materials such as paper [5], glass [6], industrial waste [7], hazardous materials [8], cartridge [9] and vegetal oil [10]. The first publication about waste collection context was made in 1974 by Beltrami and Bodin who explored techniques to solve this kind of problem [11], since then numerous papers and investigations have been published. In the literature, there are different types of models that can be applied to solve this type of problem. In essence the collection of waste is a Vehicle Routing Problem (VRP) that consists in the assignment of routes to a set of vehicles to collect waste of the clients under certain constraints [12].

Among the type of models, there are: TSP (Travelling Salesman Problem) [13], CPP (Chinese Postman Problem) [14], CARP (Capacitated Arc Routing Problem) [14], PVRP (Periodic Vehicle Routing Problem) [13], VRPTW (Vehicle Routing Problem with Time Windows) [15], CVRP (Capacitated Vehicle Routing Problem) [7] [8], Day Assignment Problem, Vehicle Assignment Problem, Day and Vehicle Assignment Problem [13], Location – Routing [16], Closed-Loop Supply Chain [9], Multi-Depot Vehicle Routing Problem with Intermediate Facilities [17].

In terms of objective functions, the most common in the literature is the minimization of costs [7], minimize the total cost of shipping recyclables from zones to centers [18], minimize the travel cost [19], minimize the operation costs [13], minimize the time and the distance [17], minimize the costs of the routes and the cost of the use of a vehicle [15], but there are also objectives such as: maximize the recuperation of waste at source [5], minimize the number of vehicles or resources to use [3], minimize distance [1], minimize the risk of materials transportation [20], minimize the time of the routes [14], maximize the waste collection [21], maximize social and environmental profits [22] and maximize the compactness of the route [3].

Although there are numerous types of models for routing in waste management systems, most of them are NP-Hard which implies that its solution cannot be reached in a polynomial time. Around combinatorial models like those mentioned before there have been done several works on methods that could reduce the computational time with acceptable solution [23].

Those methods can be classified as exact methods, heuristics and metaheuristics. For exact methods, there are application of Branch and Bound, Dynamic programming and Column generation [24]. The principal limitation of these methods is its non-polynomial computational times to find solution.

Heuristics and Metaheuristics have less computational time for solving the problem. Heuristics are specific algorithm for a problem; those try to find good solutions, not necessary the optimal one. The principal types of heuristics are: Construction heuristic, that sequentially or parallel inserts new customers into the solution; two phases heuristics, these can be Cluster first-Route second or Route first- Cluster second; and local search heuristics that from an initial solution try to improve it looking in a neighborhood of solutions [23]. The principal limitation of heuristic procedures is that the solution could be a local optimum instead of a solution close to the optimal one.

Metaheuristics have the capacity of avoid local optimums as they have better exploration in solutions space. These are generic procedures, based on analogies of nature and human process. Some of the most common types of metaheuristics are: simulated annealing, tabu search, variable neighborhood search, large neighborhood search, evolutionary algorithms and ant colonies [25], [26].

We propose a hybrid approach of mixed integer programming model and clustering to solve the selective collection service scheduling and routing problem of domestic and recyclable solid waste in the Bogotá's context. We remark that this is an extended version of a short paper recently published in the Workshop on Engineering Applications - WEA2015- that was held in Bogotá, September 2015 (see [27]).

This paper is organized as follows: Section 2 shows the problem statement. In section 3, a MIP approach is stated. Section 4 presents the methods of solution of the model. Section 5 presents examples based on the information of a work zone of two organizations in an area of Bogotá, one of them is solve by the implementation of three different cluster methods. Finally, section 6 establish some conclusions and future work lines.

## 2. Problem statement

The current solid waste management in Bogotá has into account three different programs: The Master Plan for Solid Waste Integral Management (PMIRS), Bogotá Zero Waste and the Inclusion Plan for informal recyclers [28]. The first one is the PMIRS, which was established in 2006 in order to have a global view of general waste management where territorial-environmental, social-productive and economic-financial aspects have been considered. From PMIRS there was made a pilot program in which macro-routes were implemented in order to pick recyclable waste. Those routes were the same as ordinary collection routes of the principal waste operators in the city. With this strategy, 37% of the population was covered. Despite having had this cover and having influenced the segregation waste habit of this population, there was no evidence of social inclusion [29] which was a constitutional mandate. With this premise, in 2012 the program Bogotá Zero Waste and the Inclusion Plan for informal recyclers were presented.

Bogotá Zero Waste has six principal subjects: sustainable production strategy, culture of reduction and separation at source, recyclable model for Bogotá, final use and minimization of waste disposition in Doña Juana Landfill, zero debris and integral management of special and hazardous waste.

In the Inclusion Plan for informal recyclers, they have been organized into legal organizations called ORA (Organization of Authorized Recyclers), regulated by the Administrative Special Unit of Public Services (UAESP). The principal objective is to bring them better conditions so fundamental rights could be assure in order to reduce their level of vulnerability [30].

Each ORA has a number of partners who have the recycling labor in Bogotá. They make the collection following historical routes but they do not have any technical criterion in their normal itineraries so it cannot be guaranteed the maximization of recyclable waste collection.

In this context it is proposed an optimization model for de design of selective routes, based on some assumptions which are listed below:

- Based on the program Bogotá Zero Waste, people in the city made separation at source. In this separation, ordinary waste is disposed in black bag and the recyclable one in white bag.
- Collection is made just for recyclable solid waste generated by residential users.
- The waste volume will not exceed the vehicle volume capacity. The constraint is the weight capacity.
- Users dispose their waste just one time per day.
- Ordinary waste operators make their route every two days, when this happens they collect every kind of waste so selective collection will be limited by this situation in two days. There are two consequences causes by this event; first, if recyclers do not recollect waste they will not be able to do so, as the total amount is collected by waste operators. Second, as all residues are collected the next day there is not accumulation and blocks have not waste at the beginning of the day.
- The ORA has eight-hour workday according to national legislation.
- When vehicles enter the area of collection have full capacity is available.

## 3. MIP Approach

We propose a mixed integer programming model for the collection services of recyclable waste in the Bogotá's context. This is based on the Vehicle Routing Problem (VRP) which is a generalization of TSP proposed by Dantzig and Ramser in 1959 [23]. In addition, there were made some modifications according to the real features of waste collection system in a city. In the next four sections, we describe the index sets, parameters, decision variables and mathematical formulation of the model.

### 3.1. Index sets

The index set are described as follows:

- $I : \{I_1, I_2, I_3, \dots, I_\alpha\}$  Set of blocks indexed by  $i$ .
- $J : \{J_0, J_1, J_2, J_3, \dots, J_{Deposit}\}$  Set of corners indexed by  $j$ .
- $J' : \{J_1, J_2, J_3, \dots, J_n\}$  Set of corners indexed by  $j$  from users of the service.
- $K : \{K_1, K_2, K_3, \dots, K_\sigma\}$  Set of vehicles indexed by  $k$ .
- $T : \{T_1, T_2\}$  Set of days indexed by  $t$ .
- $S : \{S_1, S_2\}$  Set of accumulation states indexed by  $s$ .

### 3.2. Parameters

The parameters are described as follows:

- $R_{jts}$ : Potential recyclable waste, dispose at corner  $j$  in the day  $t$ , with an accumulation state  $s$ ,  $j \in J'$ ,  $t \in T$ ,  $s \in S$ .
- $C_k$ : Capacity of the vehicle  $k$ ,  $k \in K$ .
- $v_k$ : Speed of the vehicle  $k$ ,  $k \in K$ .
- $T_{js}$ : Service time in corner  $j$  with an accumulation state  $s$ ,  $j \in J'$ ,  $s \in S$ .
- $d_{pq}$ : Distance of the  $arc(p, q), (p, q) \in J$ .
- $G$ : Hours of working day.
- $N$ : Number of blocks including origin and depot.
- $a_{ij}$ : 1 when the corner  $j$  belongs to the block  $i$ , 0 otherwise,  $i \in I$ ,  $j \in J$ .

### 3.3. Decision variables

The decision variables are described as follows:

- $\chi_{pqkt}$ : Take the value of 1 if an  $arc(p, q)$  is visit by the vehicle  $k$  in the day  $t$  and the value of 0 otherwise,  $p, q \in J | p \neq q, k \in K, t \in T$ .
- $w_{qkts}$ : Take the value of 1 if waste in the corner  $q$  are recollected by the vehicle  $k$ , the day  $t$  with an accumulation state  $s$  and the value of 0 otherwise,  $q \in J, k \in K, t \in T, s \in S$ .
- $u_{jkt}$ : Auxiliary variable associated to the corner  $j$  visited by the vehicle  $k$  in the day  $t$ ,  $j \in J, k \in K, t \in T$ .

### 3.4. Model

The mathematical formulation of our MIP model is described as follows:

$$MaxZ = \sum_{q \in J'} \sum_{k \in K} \sum_{t \in T} \sum_{s \in S} R_{qts} w_{qkts} \quad (1)$$

Subject to:

$$\sum_{p \in J} \sum_{q \in J'} \sum_{k \in K} \sum_{t \in T} a_{iq} \chi_{pqkt} \geq 1 \quad \forall i \in I \quad (2)$$

$$\sum_{p \in J} \sum_{q \in J'} \sum_{k \in K} a_{iq} \chi_{pqkt} \leq 1 \quad \forall i \in I, t \in T \quad (3)$$

$$\sum_{p \in J/p \neq h} \chi_{phkt} - \sum_{q \in J/q \neq h} \chi_{hqkt} = 0 \quad \forall h \in J', k \in K, t \in T \quad (4)$$

$$\sum_{q \in J/q \neq J_0} \chi_{pqkt} = 1 \quad \forall p \in J/p = J_0, k \in K, t \in T \quad (5)$$

$$\sum_{p \in J/p \neq J_{Deposit}} \chi_{pqkt} = 1 \quad \forall q \in J/q = J_{Deposit}, k \in K, t \in T \quad (6)$$

$$\sum_{p \in J'} \sum_{q \in J} \sum_{s \in S} \left( \frac{d_{pq}}{v_k} + T_{qs} \right) (\chi_{pqkt}) \leq G \quad \forall k \in K, t \in T \quad (7)$$

$$\sum_{q \in J'} \sum_{s \in S} R_{qts} w_{qkts} \leq C_k \quad \forall k \in K, t \in T \quad (8)$$

$$u_{pkt} - u_{qkt} + N \chi_{pqkt} \leq N - 1 \quad \forall p \in J/p \neq J_{Deposit}, q \in J/q \neq J_0 \wedge p \neq q, k \in K, t \in T \quad (9)$$

$$\sum_{q \in J'} \sum_{k \in K} a_{iq} w_{qkt_1 s_1} - \sum_{p \in J} \sum_{q \in J'} \sum_{k \in K} a_{iq} \chi_{pqkt_1} = 0 \quad \forall i \in I \quad (10)$$

$$\sum_{q \in J'} \sum_{k \in K} (a_{iq} w_{qkt_2 s_1} + a_{iq} w_{qkt_2 s_2}) - \sum_{p \in J} \sum_{q \in J'} \sum_{k \in K} a_{iq} \chi_{pqkt_2} = 0 \quad \forall i \in I \quad (11)$$

$$\sum_{q \in J'} \sum_{k \in K} (a_{iq} w_{qkt_1 s_1} + a_{iq} w_{qkt_2 s_2}) = 1 \quad \forall i \in I \quad (12)$$

$$\sum_{s \in S} w_{qkts} - \sum_{p \in J} \chi_{pqkt} = 0 \quad \forall q \in J', k \in K, t \in T \quad (13)$$

$$\chi_{pqkt} \in \{0, 1\} \quad \forall p, q \in J, k \in K, t \in T \quad (14)$$

$$w_{qkts} \in \{0, 1\} \quad \forall q \in J, k \in K, t \in T, s \in S \quad (15)$$

$$u_{jkt} \geq 0 \quad \forall j \in J, k \in K, t \in T \quad (16)$$

The objective function (1) maximizes the amount of waste recollected. Constraints (2) establish that Blocks must be visited at least one time in two days. Constraints (3) limit visits per day to one. Constraints (4) are the balance flow. Constraints (5) and (6) force vehicles must start at origin, along one unique route, and finish at depot. Constraints (7) limit the workday hours according to national legislation. Constraints (8) put the maximum capacity of each vehicle. Constraints (9) eliminate sub tours. Constraints (10) ensure that when an arc is activated there is waste recollection in day one. Constraints (11), as well as equation (10), ensure recollection and also these do not allow double accumulation state in day two. Constraints (12) control the state of day one and day two to prevent accumulation in day two when collection have been done in day one. Constraints (13) show the relation between the  $\chi_{pqkt}$  arc variable and the recollection variable  $w_{qkts}$ . Constraints (14) and (15) determine the nature of variables  $\chi_{pqkt}$  and  $w_{qkts}$  as binary. Constraints (16) establish the non-negativity of variable  $u_{jkt}$ .

## 4. Solution methods

To solve the model proposed in section 3.4, we proposed a method of *cluster-first route-second* in which capacity constraint of each vehicle assigned to a cluster must be respected. The strategy is to create clusters grouping blocks and then solve the proposed model. Moreover, there are applied three different methods of cluster. Its purpose is to evaluate which one brings better solutions.

### 4.1. A Centroid-based heuristic algorithm

This algorithm is based on the geometry of geometrical centers, around which the clusters are generated. This method is divided into two phases. In the first phase, clusters are constructed by selecting the farthest node from the origin, among the nodes that have not been assigned, and it is generated a first cluster; then the geometrical center of the cluster (17), in order to add the nodes closer to the GC, considering the defined capacity of each cluster [31], [32].

$$GC(l_i) = \left( \sum_{i=0}^m \frac{w_i^x}{m}, \sum_{i=0}^m \frac{w_i^y}{m} \right) \quad (17)$$

Where  $w_i^x$  and  $w_i^y$  are the coordinates in x, y of the nodes that belong to the cluster. Table I shows the pseudocode of the first phase. After of generating clusters, it is necessary to adjust them in a second phase, in Table II the process is shown.

**Table I.** Pseudo code for the Centroid-based Heuristic Algorithm, Phase I

---

```

Start phase I
Set  $i = 0$  and  $Q = \text{vehicle capacity}$ 
while (unvisited block exist)
     $v_j$  = the frathest block among un-clustered blocks from the origin
    Generate cluster  $l_i$  with  $v_j$ 
    Set capacity of  $l_i = Q$ 
    while (waste of  $v_j$  do not exceed available capacity of  $l_i$ )
        Add  $v_j$  to  $l_i$ 
        Reduce available capacity of  $l_i$  by  $v_j$ 
        Calculate  $GC(l_i)$ 
         $v_j$  = the closest block among un-clustered blocks from  $GC(l_i)$ 
    End while
     $i = i + 1$ 
End while
End phase I
    
```

---

**Table II.** Pseudo code for the Centroid-based Heuristic Algorithm, Phase II

---

```

Start phase II
Set  $L = \{l_0, l_1, \dots, l_m\}$ 
For  $i = 0$  to  $m$  repeat
    For every block  $v_k$  in  $l_i$ 
        For every  $l_j$  in  $L$ 
            If ( $i \neq j$  and  $v_k$  is closer to  $GC(l_j)$  than  $GC(l_i)$  and available capacity of  $l_j$ 
            is equal to or bigger than waste of  $v_k$ )
                Move  $v_k$  from  $l_i$  to  $l_j$ 
                Recalculate  $GC(l_i)$  and  $GC(l_j)$ 
            End if
        End for
    End for
End phase II
    
```

---

## 4.2. Sweep algorithm

This heuristic gives shape to the cluster based on the geometry of polar coordinates. From an origin point, there is displayed a straight that rotates on the zone where nodes must be assigned. The area cover by this straight is a cluster only while capacity constraints are respected [33]. Table III shows the pseudocode.

**Table III.** Pseudo code for the of the Sweep Algorithm

---

```

Start initialization
Set  $i = 0$  and  $Q = \text{Vehicle capacity}$ 
Calculate polar coordinates of blocks  $(\rho_i, \theta_i)$ 
Sort the blocks by  $\theta_i$ 
If ( $\theta_i = \theta_{i+1}$ )
    Sort the blocks by  $\rho_i$ 
End If
End Initialization
Start cluster generation
Select an initial block  $v_j$ 
Do  $k = 1$  y  $C_k = \{v_j\}$ 
while (unvisited block exist)
    Visit the next block  $v_{j+1}$ 
    If (Waste generated in the block  $v_{j+1}$  do no break capacity restrictions)
         $C_k = C_k \cup \{v_{j+1}\}$ .
    Else
         $k = k + 1$  and create  $C_k = \{v_j\}$ 
    End if
End while
End cluster generation
    
```

---

### 4.3. Proposed localization model

Another way to establish cluster is the application of localization models. Fisher and Jaikumar proposed a Generalized Assignment Problem (GAP) for which initial seeds must be selected [34]. Another approach was proposed by Bramel and Simchi-Levi in which some nodes are previously selected to be possible locations for concentrators [33]. Around these concentrators other nodes are assigned as terminals and then, cluster are determined; for this model seeds are found from the solution of the problem, but there must be previous selection of possible locations [33]. We propose a modified approach in which concentrators of cluster can be selected from the entire set of nodes.

#### 4.3.1. Index sets

The index set are described as follows:

- $I : \{I_1, I_2, I_3, \dots, I_m\}$  Set of blocks indexed by  $i$ .
- $J : \{J_0, J_1, J_2, J_3, \dots, J_n\}$  Set of concentrators indexed by  $j$ .
- $K : \{K_1, K_2, K_3, \dots, K_\sigma\}$  Set of vehicles indexed by  $k$ .

#### 4.3.2. Parameters

The parameters are described as follows:

- $f_i$ : Distance between the origin and a concentrator  $j$ .
- $c_{ij}$ : Distance between a concentrator  $j$  and a block  $i$ .
- $R_i$ : Amount of waste generated by the block  $i$ .
- $C_{jk}$ : Capacity of the concentrator  $j$  for the assigned vehicle  $k$ .

#### 4.3.3. Decision variables

The decision variables are described as follows:

- $x_{ij}$ : Take the value of 1 if block  $i$  is assigned to the concentrator  $j$ , 0 otherwise.
- $y_{jk}$ : Take the value of 1 if the vehicle  $k$  is assigned to the concentrator  $j$ , 0 otherwise.

#### 4.3.4. Model

The mathematical formulation of our MIP model is described as follows:

$$\text{Min } Z = \sum_{j \in J} \sum_{k \in K} f_j y_{jk} + \sum_{i \in I} \sum_{j \in J} c_{ij} x_{ij} \quad (18)$$

Subject to:

$$\sum_{i \in I} x_{ij} = 1 \quad \forall i \in I \quad (19)$$

$$\sum_{k \in K} C_{jk} y_{jk} - \sum_{i \in I \setminus \{0\}} R_i x_{ij} \geq 0 \quad \forall j \in J \quad (20)$$

$$\sum_{j \in J} y_{jk} = 1 \quad \forall k \in K \quad (21)$$

$$x_{ij} \in \{0, 1\} \quad \forall i \in I, j \in J \quad (22)$$

$$y_{jk} \in \{0, 1\} \quad \forall j \in J, k \in K \quad (23)$$

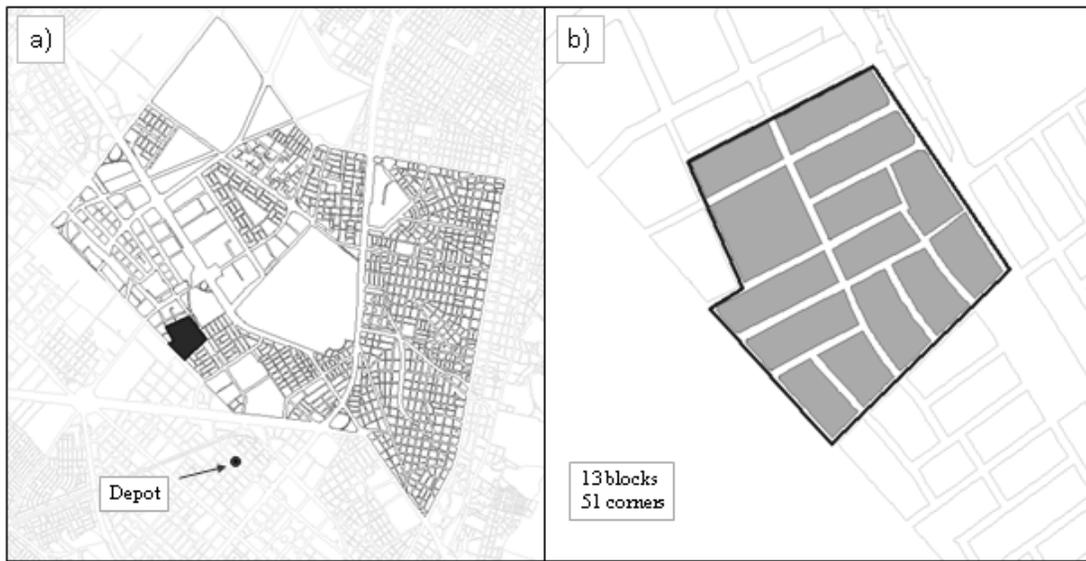
The objective function (18) minimizes the distance between the origin and each block that could be a concentrator and the distance between concentrators and the blocks associated with a concentrator. Constraints (19) force each block to belong only to one concentrator. Constraints (20) ensure that cluster capacity is respected. Constraints (21) limit to one the number of concentrators assigned to each vehicle. Constraints (22) and (23) define the nature of variables  $x_{ij}$  and  $y_{jk}$  as binaries.

## 5. Results

We run some experiments in the context of waste collection system in Bogotá. We selected two ORAs that do collection in the area of Teusaquillo, where there are eight organizations that have a delimited zone of work according to their historical routes. The ORAs selected are “La Colombianita” (see Figure 1) whose area has less users and it is possible to design a route with the proposed model in an acceptable computational time. In addition, we selected the ORA called “EMRS” (see Figure 2) that has more blocks in order to apply cluster methods and generate a comparison.

### 5.1. Case 1: La Colombianita

Table IV and Table V presents the information about “La Colombianita” case.



**Figure 1.** (a) Work zone of “La Colombianita” in Teusaquillo, area of Bogotá. (b) Blocks and corners of the work zone of “La Colombianita”.

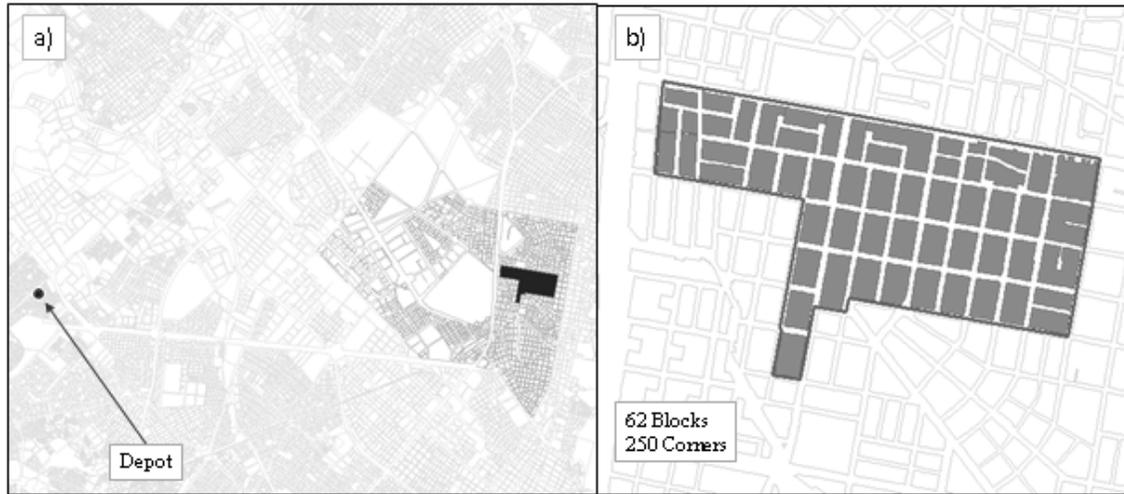
Users	3.860
Number of houses	1.453
Number of blocks	13
Number of corners	51
Amount of potentially recyclable waste per day (Kg)	302,2187

Type	Quantity	Speed	Capacity
Truck	2	23	2.800
Wheelbarrow	1	4,4	120

Trucks were codified as K1 and K2, and the wheelbarrow was codified as K3. Corners were organized by a number from 1 to 51. The distances of combination among corners were calculated applying Floyd-Warshall Algorithm, considering the real direction of the streets. The legal workday is eight hours. In the solution (see Table VI)  $J_0$  corresponds at origin and  $J_{52}$  is the depot. Each corner belongs to a different block and every constraint is satisfied.

Vehicle – Day	Route	Collected Waste (Kg)	Distance (m)
K1 – T1	J0-J32-J26-J41-J39-J52	32,41413	669,0741159
K2 – T1	J0-J22-J7-J29-J2-J47-J15-J11 -J52	236,372605	1.937,41542
K1 – T2	J0-J10-J30-J39-J47-J20-J52	161,365995	1.490,099395
K2 – T2	J0-J33-J15-J4-J8-J26-J43-J49-J16 -J52	174,28467	2.164,435921
<b>Total</b>		<b>604,4374</b>	<b>6.891,190868</b>

## 5.2. Case 2: EMRS



**Figure 2.** (a) Work zone of “EMRS” in Teusaquillo, area of Bogotá. (b) Blocks and corners of the work zone of “EMRS”.

In Table VII and Table VIII is presented the information about “EMRS”.

**Table VII. Parameters of “EMRS”**

Users	8.741
Number of houses	3.796
Number of blocks	62
Number of corners	250
Amount of potentially recyclable waste per day (Kg)	651,82

**Table VIII. Mean of Transportation of “EMRS”**

Type of transport	Quantity	Speed (Km/h)	Capacity (Kg)
Truck	1	23	800

For each clustering method, we have joined the results of every cluster in order to obtain just one route. The connections were made having into account the distance between the origin, the first and last nodes of the clusters.

Table IX shows the results with the centroid-based heuristic, where 0 corresponds at origin and  $F$  is the depot. Each corner belongs to a different block and every constraint is satisfied.

**Table IX. Results with the centroid-based heuristic of “EMRS”**

Day	Route	Distance (m)	Collected Waste (Kg)
T1	0 – 133 – 147 – 190 – 136 – 39 – 246 – 130 – 145 – 56 – 119 – 106 – 138 – 115 – 111 – 69 – 75 – 218 – 224 – 23 – 211 – 244 – 8 – 12 – 43 – 19 – 248 – 27 – 155 – 11 – 40 – 163 – 31 – 206 – 195 – 194 – 51 – F	47006,47	435,88
T2	0 – 781 – 873 – 680 – 773 – 789 – 753 – 744 – 760 – 754 – 697 – 766 – 709 – 856 – 703 – 783 – 711 – 748 – 2.688 – 898 – 993 – 331 – 665 – 389 – 693 – 330 – 388 – 852 – 3.350 – 669 – 834 – 678 – 3.305 – 738 – 659 – 700 – 685 – 687 – 735 – 875 – 770 – 827 – 880 – 870 – 774 – 730 – 902 – 714 – 988 – 726 – 857 – 861 – 695 – 890 – 883 – 692 – 872 – 878 – F	67174,06	782,43
Total		114.180,51	1.218,31

Table X presents the results with the sweep algorithm, where 0 corresponds at origin and  $F$  is the depot. Each corner belongs to a different block and every constraint is satisfied.

**Table X.** Results with the sweep heuristic of “EMRS”

Day	Route	Distance (m)	Collected Waste (Kg)
T1	0-3.350 – 709 – 682 – 874 – 783 – 780 – 786 – 678 – 902-2.688 – 899 – 1.054 – 741 – 748 – 756 – 753 – 735 – 884 – 872 – 891 – 877 – F	51.458,70	391,71
T2	0-769 – 712 – 760 – 894-1.299 – 658 – 904 – 899 – 857-3.324 – 825 – 868 – 684 – 671 – 848-1.054 – 750 – 746 – 736 – 755 – 741 – 701 – 720 – 777 – 887 – 679 – 780 – 786 – 689 – 783 – 873 – 988 – 906 – 707 – 662 – 668 – 675 – 696 – 691 – 879 – 864 – 331 – 391 – 327 – 992 – 655 – 722 – 829 – 729 – 994 – 730 – 854– 835 – 865 – 988 – 906 – 707 – 662 – 668 – 675 – F	63.380,90	796,77
Total		114.839,60	1.188,48

Table XI shows the results with the sweep algorithm, where 0 corresponds at origin and F is the depot. Each corner belongs to a different block and every constraint is satisfied.

**Table XI.** Results with the location model proposed of “EMRS”

Day	Route	Distance (m)	Collected Waste (Kg)
T1	0 – 758 – 789 – 694 – 751 – 754 – 781 – 769 – 745 – 761 – 755 – 692 – 777 – 750 – 780 – 687 – 730 – 678 – 725 – 906 – 850 – 726 – 677 – 661 – 331 – 668 – 662 – 387 – 708– 2.688 – 719 – 654 – 716 – 713 – 859 – 705 – 332– 3.305 – 832– 1.052 – 836 – 985 – 672 – 989 – 739– 1.010 – 846 – 827 – 735 – 995 – 701 – 697 – F	46.921,29	456,89
T2	0 – 770 – 744 – 687 – 690 – 752 – 779 – 3.351 – 1.053 – 832 – 987 – 835 – 682 – 845 – 738 – 1.009 – 996– 734 – 992 – 825 – 332 – 705 – 747 – 686 – 754 – 719 – 390 – 717 – 712– 2.688 – 858 – 906 – 849 – 722 – 733 – 677 – 679 – 853 – 727 – 388 – 709 – 663 – 667 – 331 – 659 – 769 – 758 – 694 – 699 – F	49.335,03	548,91
Total		96.256,32	1.005,80

For EMRS, the potential amount of waste to be collected was 1.303,64 Kg in both days. With the centroid-base heuristic were collected 91%, with sweep algorithm 93% and with localization model 77%. In a first view of the results, we could infer that the best solution method is the sweep algorithm. Nevertheless, there must be accomplished several tests to have a strong validation of the performance of each method.

## 6. Conclusions

We present a MIP approach model for selective waste collection as a variant of a VRP model. The problem was applied in real world case in an area of Bogotá, Colombia. From literature review, in selective waste collection there are several models and solution approach. In general, models depend on waste collection necessities and its environment. The selected MIP approach was VRP that takes into account additional aspects such as corners which are the nodes that must belong to specific blocks. Waste is generated by each block and its accumulation is allowed just until ordinary collection occurs.

In the solution with the ORA “La Colombianita”, all recyclable waste are collected so it can be ensured 100% of its return to the productive cycle. The model proposed can be applied in contexts where separation at source is a regular activity, and also where there are fixed points so that users could dispose their waste.

We analyzed zones with more number of blocks that are served by other organizations. However, it could not be found an optimal solution in a reasonable computational time. So we perform some test with three clustering methods. These methods could have influenced the solution when the MIP approach was applied, since all off them look for the minimization of distance within the cluster. For these reasons, initial parameters have some improvement in the reduction of distances which could mean more efficiency but not necessarily higher coverage. Future work will concentrate on validation of results in order select the most appropriate method.

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